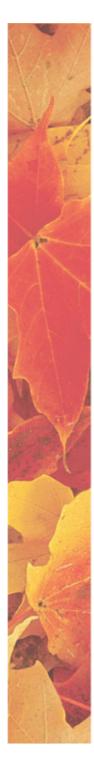


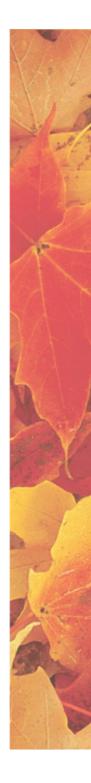
Recursion

 Recursion is a fundamental programming technique that can provide an elegant solution certain kinds of problems



Recursive Thinking

- A recursive definition is one which uses the word or concept being defined in the definition itself
- When defining an English word, a recursive definition is often not helpful
- But in other situations, a recursive definition can be an appropriate way to express a concept
- Before applying recursion to programming, it is best to practice thinking recursively



Recursive Definitions

Consider the following list of numbers:

24, 88, 40, 37

• Such a list can be defined as follows:

A LIST is a: number or a: number comma LIST

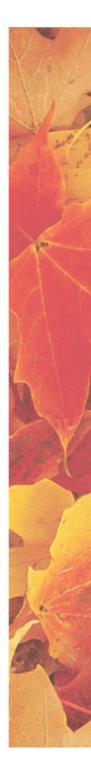
- That is, a LIST is defined to be a single number, or a number followed by a comma followed by a LIST
- The concept of a LIST is used to define itself



Recursive Definitions

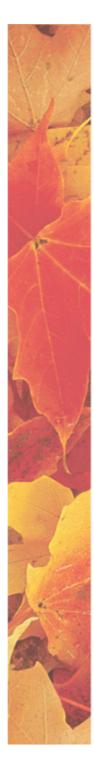
• The recursive part of the LIST definition is used several times, terminating with the non-recursive part:

number	comma	LIST					
24	,	88, 40, 37					
		number 88	comma				
				numk	er	comma	LIST
				40)	1	37
							number 37



Infinite Recursion

- All recursive definitions have to have a nonrecursive part
- If they didn't, there would be no way to terminate the recursive path
- Such a definition would cause *infinite recursion*
- This problem is similar to an infinite loop, but the non-terminating "loop" is part of the definition itself
- The non-recursive part is often called the base case

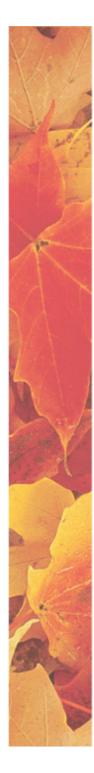


Recursive Definitions

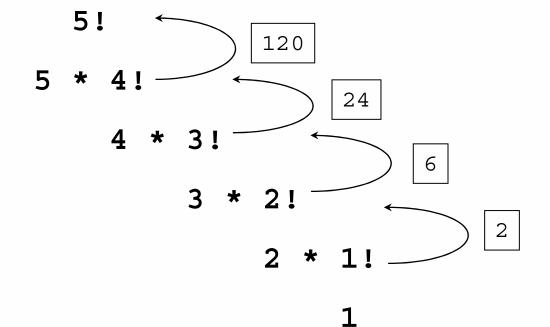
- N!, for any positive integer N, is defined to be the product of all integers between 1 and N inclusive
- This definition can be expressed recursively as:

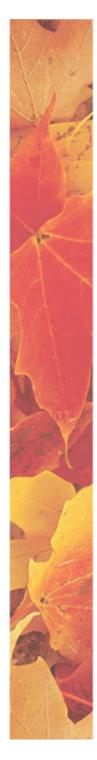
$$N! = N * (N-1)!$$

- A factorial is defined in terms of another factorial
- Eventually, the base case of 1! is reached

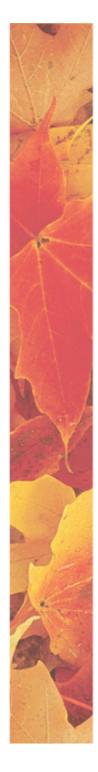


Recursive Definitions





- A Function can invoke itself; if set up that way, it is called a *recursive function*
- The code of a recursive function must be structured to handle both the base case and the recursive case
- As with any function call, when the function completes, control returns to the function that invoked it (which may be an earlier invocation of itself)



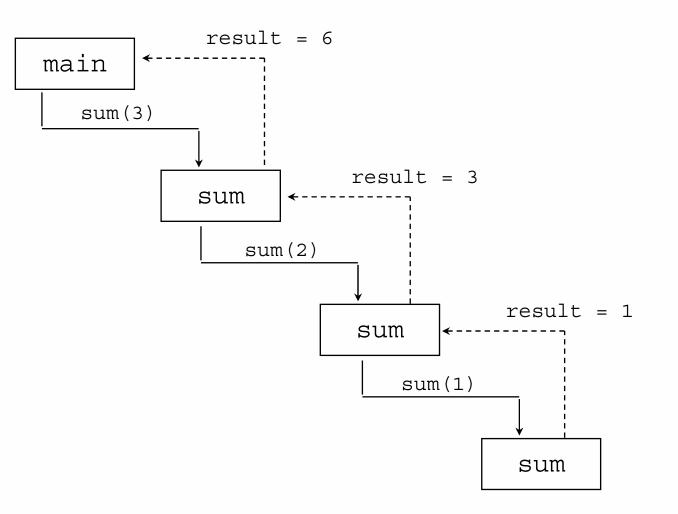
- Consider the problem of computing the sum of all the numbers between 1 and any positive integer N
- This problem can be recursively defined as:

$$\sum_{i=1}^{N} i = N + \sum_{i=1}^{N-1} i = N + N-1 + \sum_{i=1}^{N-2} i$$
$$= N + N-1 + N-2 + \sum_{i=1}^{N-3} i$$



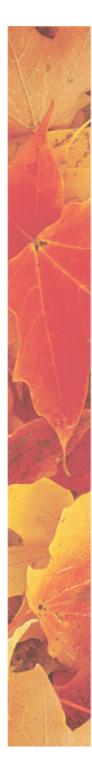
```
// This function returns the sum of 1 to num
int sum (int num)
{
    int result;
    if (num == 1)
        result = 1;
    else
        result = num + sum (n-1);
    return result;
}
```





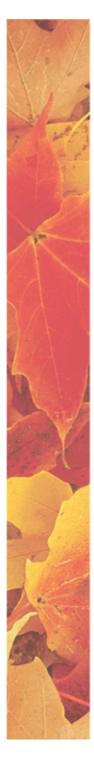


- Note that just because we can use recursion to solve a problem, doesn't mean we should
- For instance, we usually would not use recursion to solve the sum of 1 to N problem, because the iterative version is easier to understand
- However, for some problems, recursion provides an elegant solution, often cleaner than an iterative version
- You must carefully decide whether recursion is the correct technique for any problem

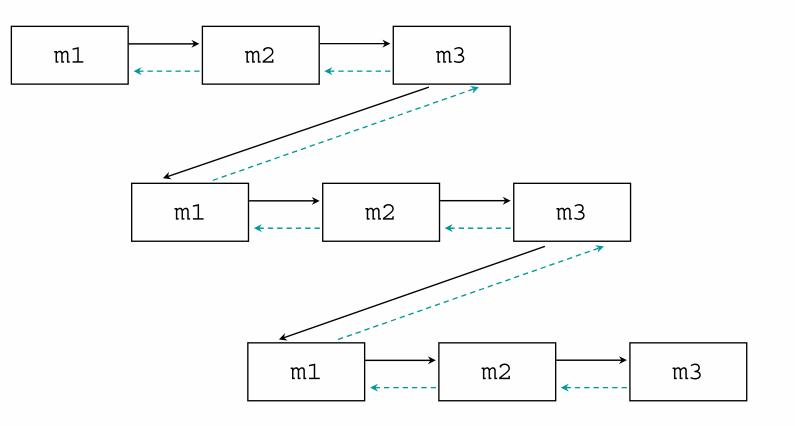


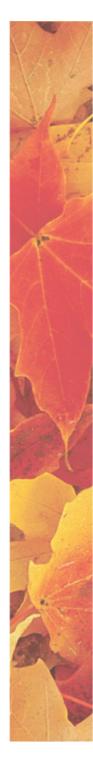
Indirect Recursion

- A function invoking itself is considered to be *direct recursion*
- A function could invoke another function, which invokes another, etc., until eventually the original function is invoked again
- For example, function m1 could invoke m2, which invokes m3, which in turn invokes m1 again
- This is called *indirect recursion*, and requires all the same care as direct recursion
- It is often more difficult to trace and debug



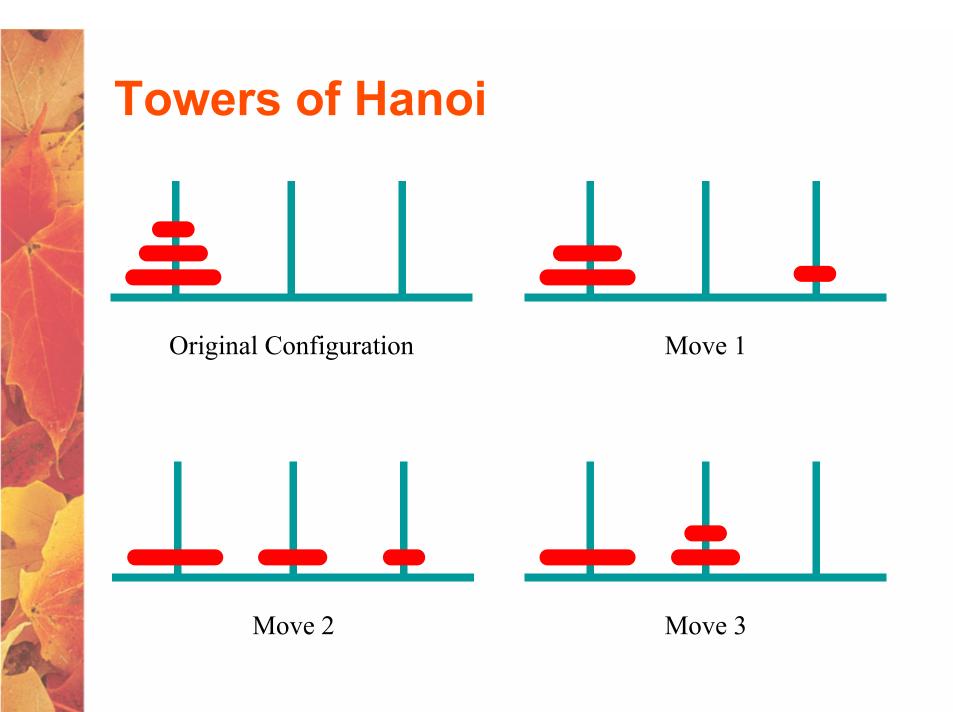
Indirect Recursion

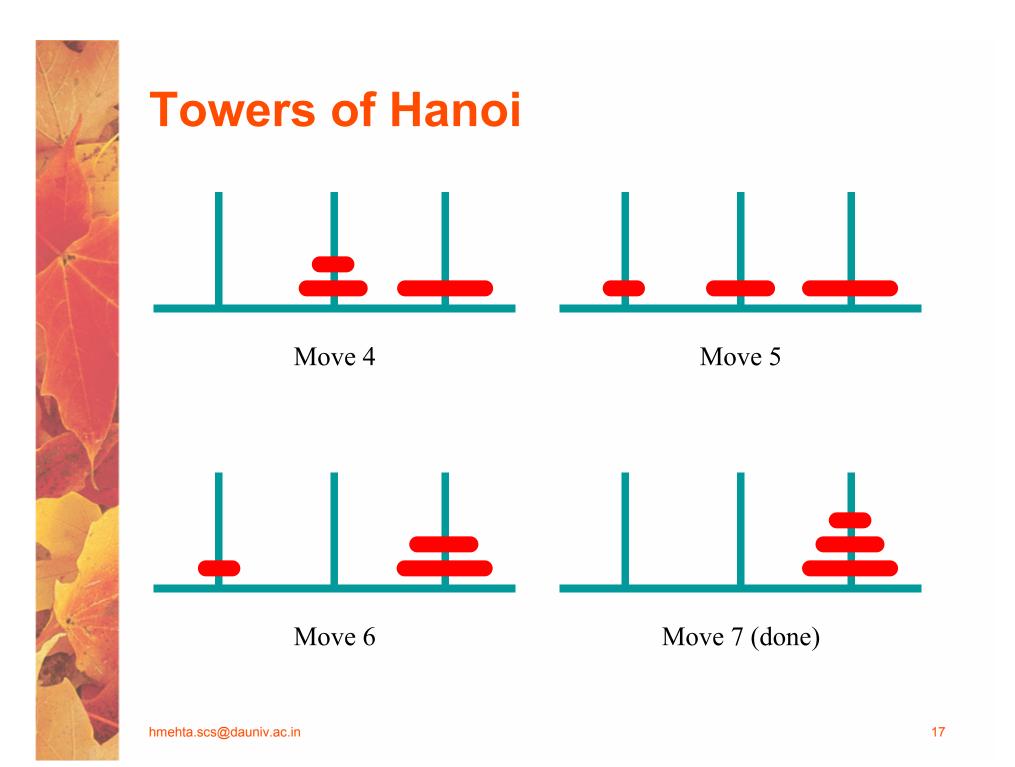




Towers of Hanoi

- The Towers of Hanoi is a puzzle made up of three vertical pegs and several disks that slide on the pegs
- The disks are of varying size, initially placed on one peg with the largest disk on the bottom with increasingly smaller ones on top
- The goal is to move all of the disks from one peg to another under the following rules:
 - We can move only one disk at a time
 - We cannot move a larger disk on top of a smaller one







Towers of Hanoi

- An iterative solution to the Towers of Hanoi is quite complex
- A recursive solution is much shorter and more elegant



Towers of Hanoi

#include <stdio.h>
#include <conio.h>

void transfer(int,char,char,char);

int main()

```
{
```

}

}

```
int n;
printf("Recursive Solution to Towe of Hanoi Problem\n");
printf("enter the number of Disks");
scanf("%d",&n);
transfer(n,'L','R','C');
getch();
return 0;
```

void transfer(int n,char from,char to,char temp)

```
if (n>0)
{
    transfer(n-1,from,temp,to); /* Move n-1 disk from origin to temporary */
    printf("Move Disk %d from %c to %c\n",n,from,to);
    transfer(n-1,temp,to,from); /* Move n-1 disk from temporary to origin */
}
return;
```



Drawbacks of Recursion

Regardless of the algorithm used, recursion has two important drawbacks:

- Function-Call Overhead
- Memory-Management Issues

Eliminating Recursion — Tail Recursion

A special kind of recursion is tail recursion.

Tail recursion is when a recursive call is the last thing a function does.

Tail recursion is important because it makes the recursion \rightarrow iteration conversion very easy.

- That is, we like tail recursion because it is easy to eliminate.
- In fact, tail recursion is such an obvious thing to optimize that some compilers automatically convert it to iteration.

Eliminating Recursion — Tail Recursion

For a void function, tail recursion looks like this:

```
void foo(TTT a, UUU b)
{
    ...
    foo(x, y);
}
```

For a function returning a value, tail recursion looks like this:

```
SSS bar(TTT a, UUU b)
{
    ...
    return bar(x, y);
}
```



A tail-recursive Factorial Function

We will use an auxiliary function to rewrite factorial as tailrecursive:

```
int factAux (int x, int result)
{
    if (x==0) return result;
    return factAux(x-1, result * x);
}
int tailRecursiveFact( int x)
{
    return factAux (n, 1);
}
```